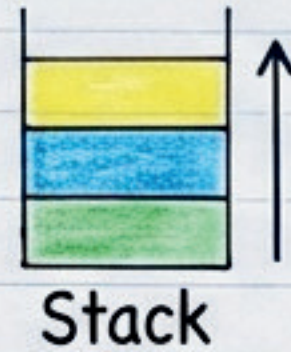
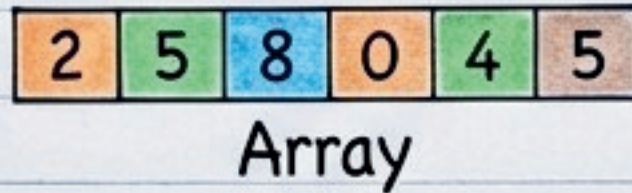


# Chapter 1: Introduction to DSA

(notes by ~ Kamal kishor)

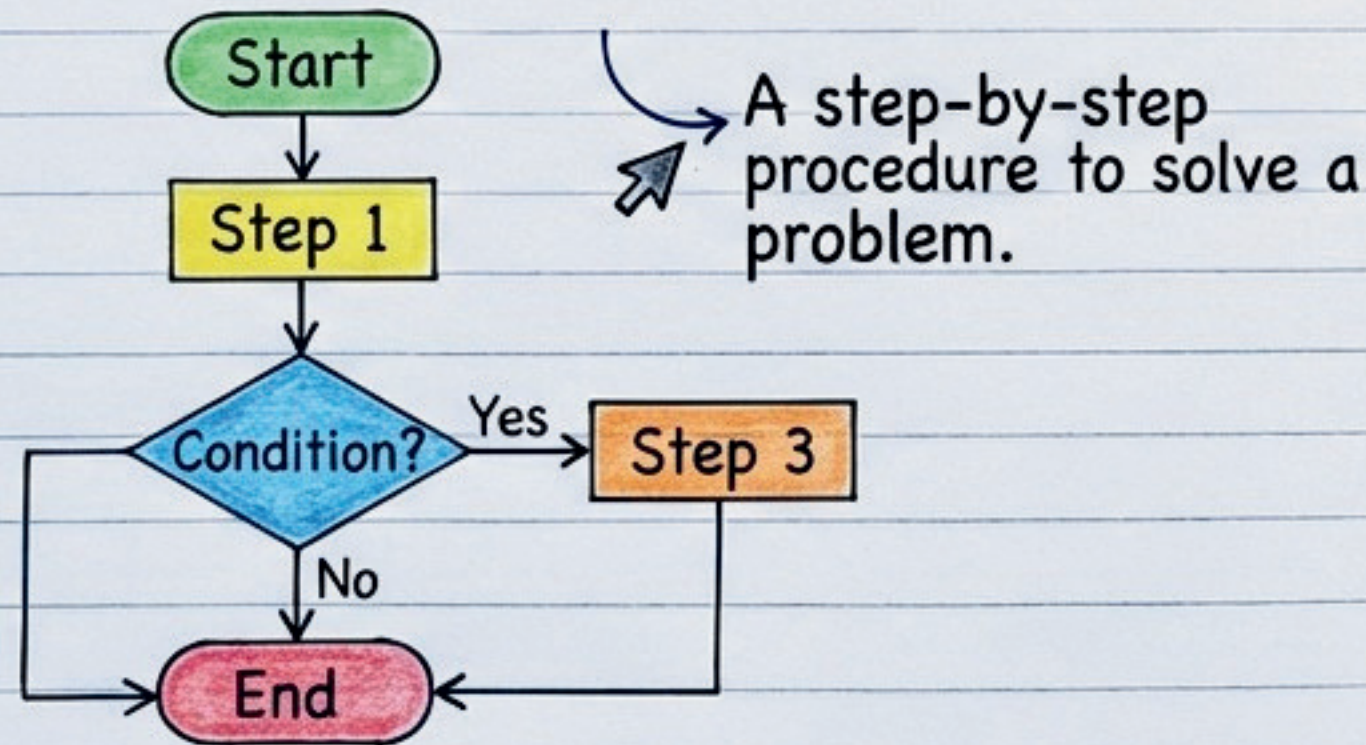
## What is Data Structure?

A way to store and organize data in a computer so that it can be used efficiently.



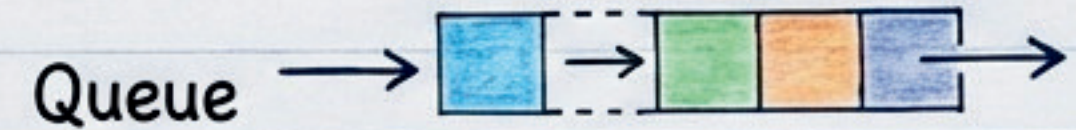
## What is Algorithm?

A finite sequence of well-defined steps to solve a specific problem. Like a recipe!



## Properties of Algorithms

- ✓ 1. Finiteness: Must terminate after a finite number of steps.
- ✓ 2. Definiteness: Steps must be precisely defined.
- ✓ 3. Input/Output: Takes inputs, produces outputs.

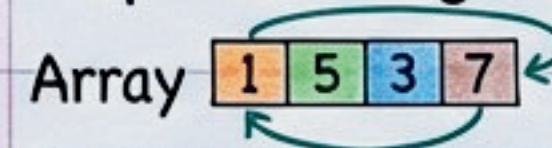


## Why DSA is important?

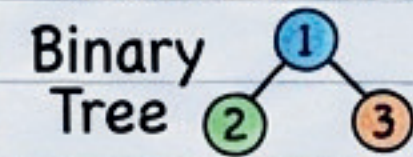
- ✓ Efficient Problem Solving (choosing right tools).
- ✓ Optimal Resource Usage (memory/time).
- ✓ Coding Interviews (crucial for challenges).

## Types of Data Structures

Linear  
Sequential organization.



Non-Linear  
Hierarchical.



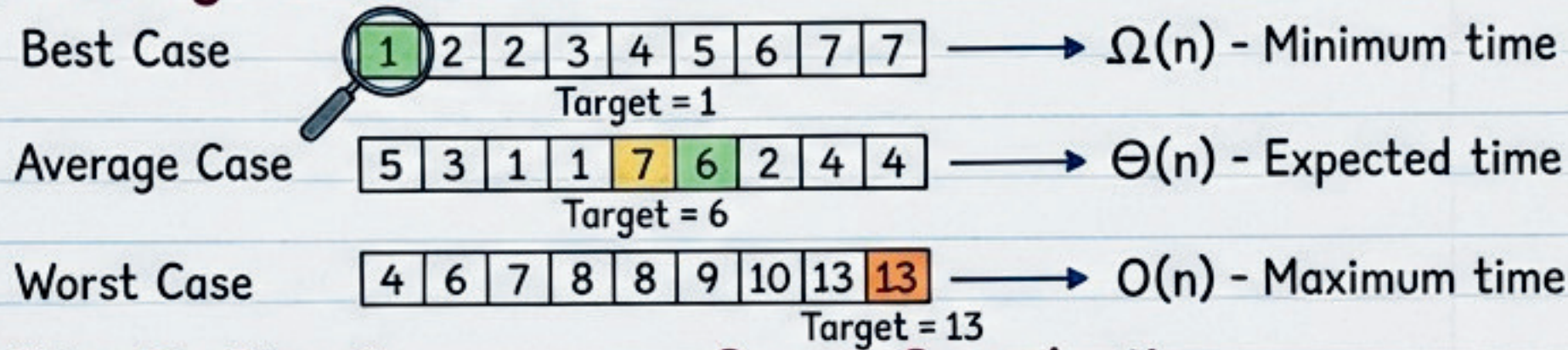
**Abstract Data Types (ADT):** Defined by operations (List, Stack, Queue, Map).

# Time & Space Complexity

## Why complexity analysis?

- ✓ Evaluate Efficiency: Measuring running time & memory usage.
- ✓ Compare Algorithms: Choosing the optimal one.
- ✓ Scalability: Handling large inputs without degrading.

## Best, Average & Worst Case

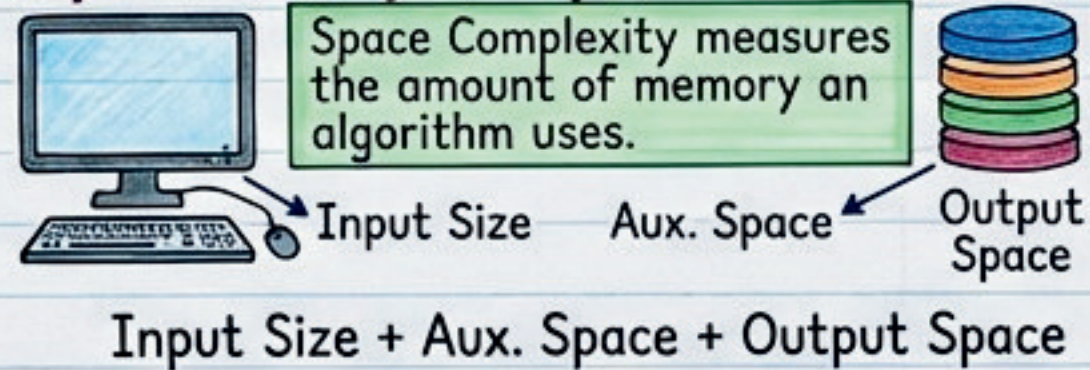


## Big-O, Big- $\Omega$ , Big- $\Theta$

1	↑	Big-O	Upper Bound	$f(n) \leq O(g(n))$
2	↓	Big- $\Omega$	Lower Bound	$f(n) \geq \Omega(g(n))$
3	↕	Big- $\Theta$	Tight Bound	$f(n) = \Theta(g(n))$

These notations provide a measure of the time complexity of an algorithm in relation to the input size.

## Space Complexity



## Complexity of Loops & Recursion

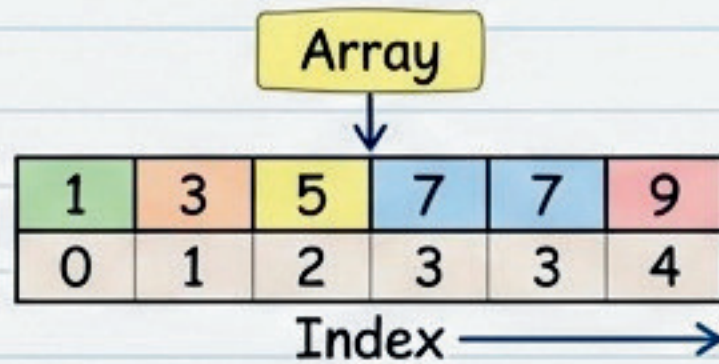
`for (int i=0; i<n; i++)`  $\rightarrow O(n)$

`for (i... for j...)`  $\rightarrow O(n^2)$

`for (i=1; i<n; i*=2)`  $\rightarrow O(\log n)$

# Chapter 2: Arrays (Introduction & Types)

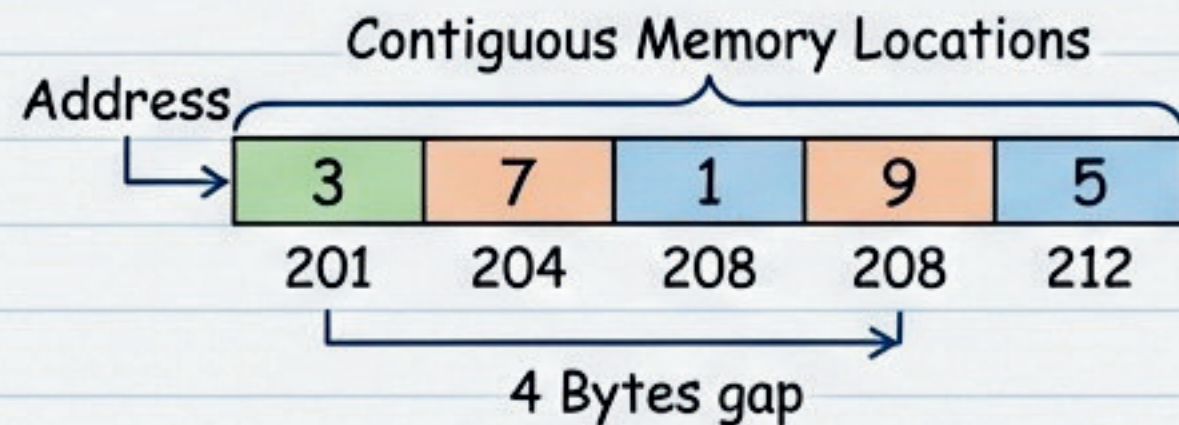
An **array** is a collection of elements of the same data type stored in contiguous memory locations.



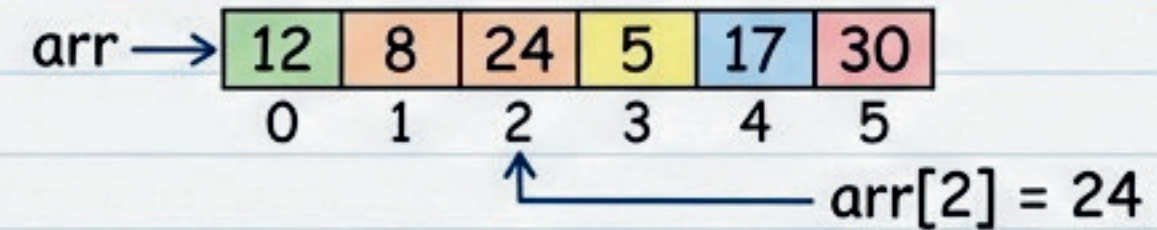
## Characteristics of Arrays

<b>Fixed Size:</b> Defined at declaration.	<b>Homogeneous Elements:</b> Same data type.	<b>Contiguous Memory:</b> Adjacent locations.
---	---	--

## Memory Representation

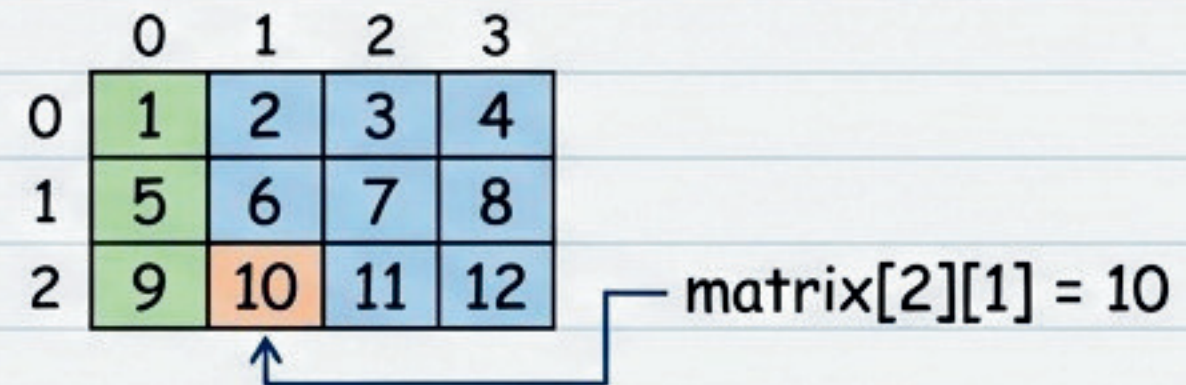


## 2.2 1D Arrays

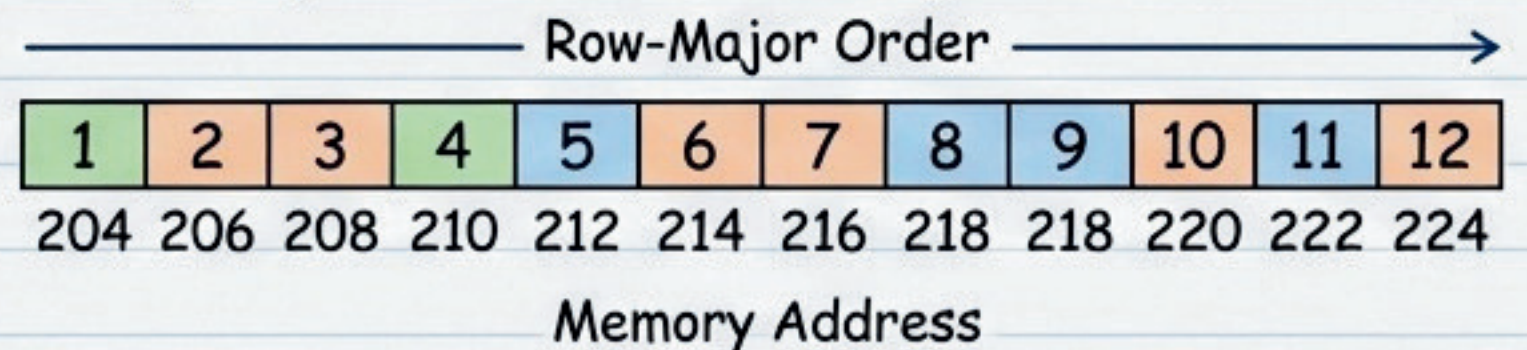


```
int arr[5] = {3, 7, 1, 9, 5};
```

## 2.3 2D Arrays (Matrix)

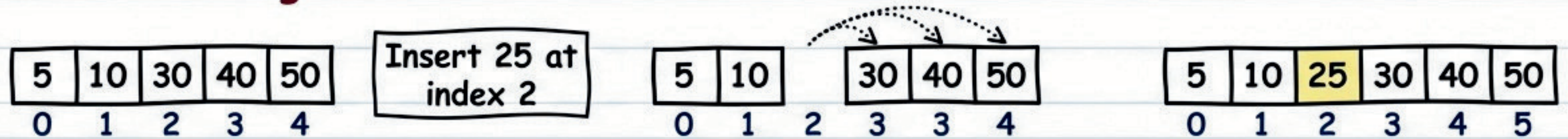


## Memory Representation (Row-Major Order)



# Array Operations & Problems

## 1. Insertion Logic



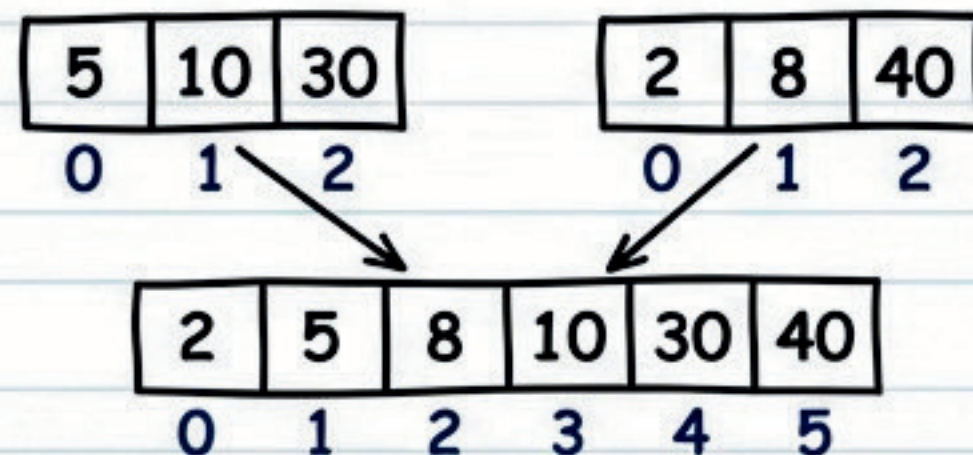
✓ Time Complexity  $\rightarrow O(n)$  (Worst case shifting)

## 2. Deletion Logic



✓ Time Complexity  $\rightarrow O(n)$

## 3. Merge Two Sorted Arrays



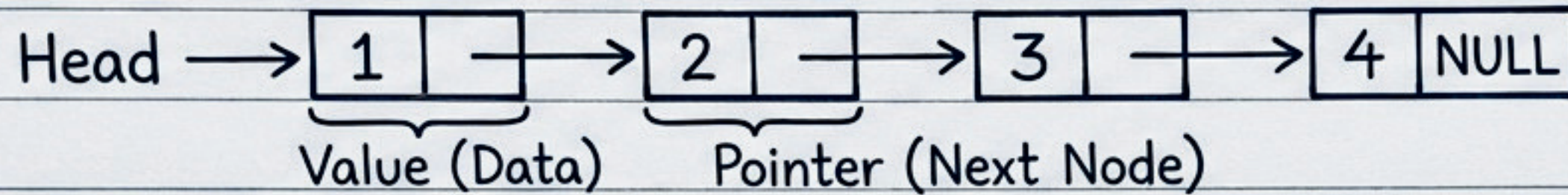
### Tips for Solving Array Problems

- ✓ Plan approach before coding.
- ✓ Consider edge cases (empty array, large values).
- ✓ Watch out for fragmentation.

# Chapter 3: Linked List (Basics)

## 3.1 Introduction

A sequence of elements (nodes) where each contains a value and a pointer to the next.



## 3.2 Singly Linked List

✓ Unidirectional

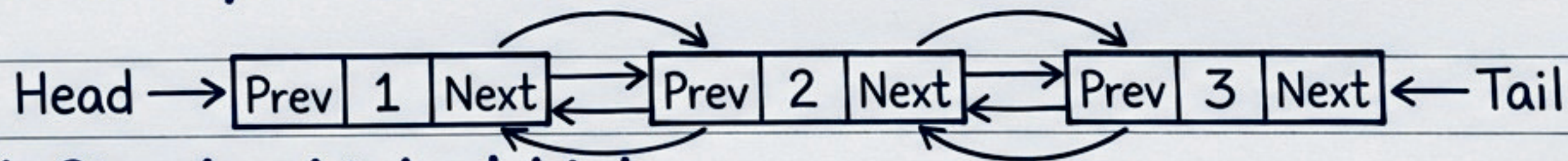
✓ Dynamic Size

✓ Memory Efficient

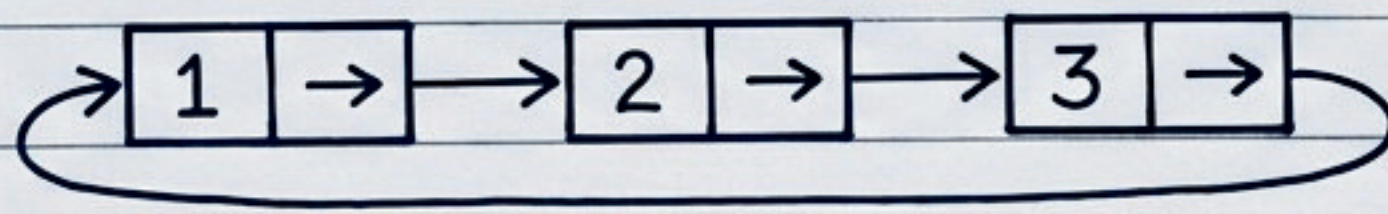
How to create a new node? ● New Node → 

10	Next
----	------

## 3.3 Doubly Linked List



## 3.4 Circular Linked List

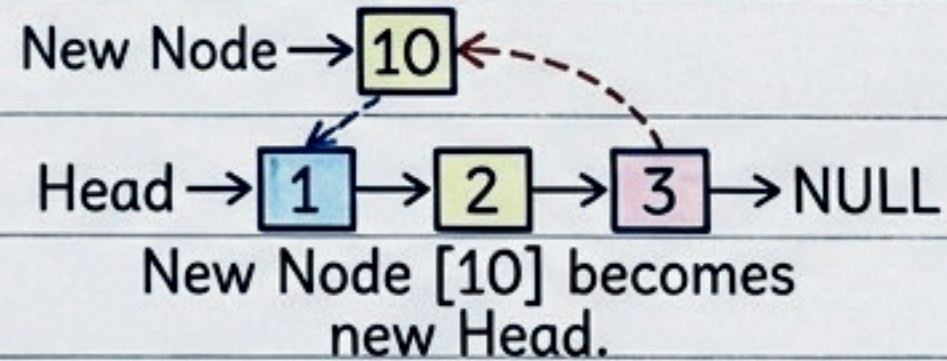


No NULL value; forms a continuous loop.

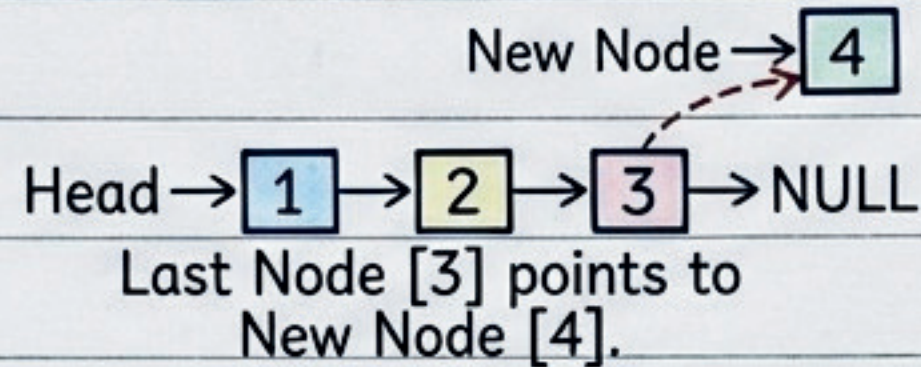
# Linked List Operations

## 1. Insertion Logic

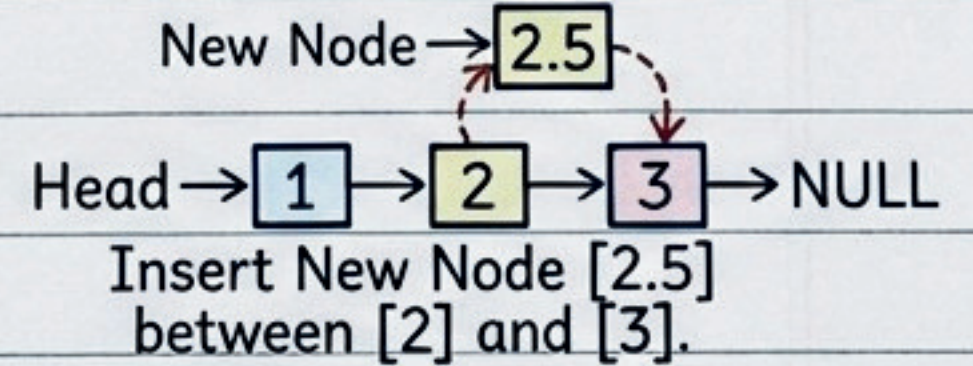
### Beginning



### End

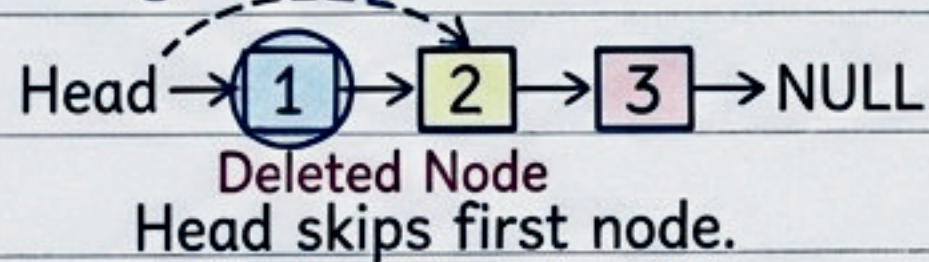


### Specific Position

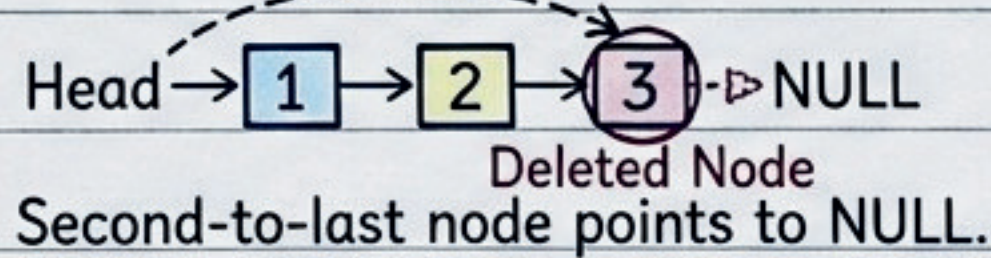


## 2. Deletion Logic

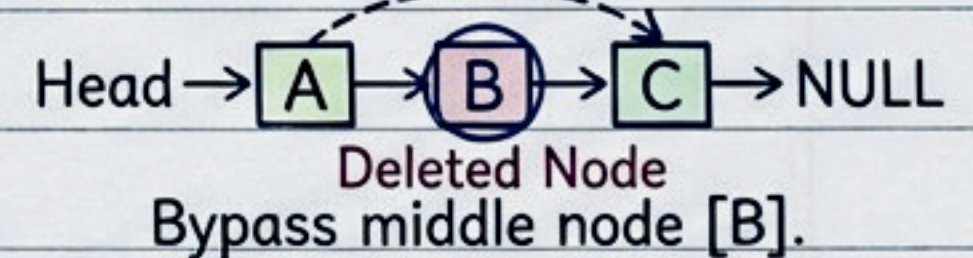
### Beginning



### End



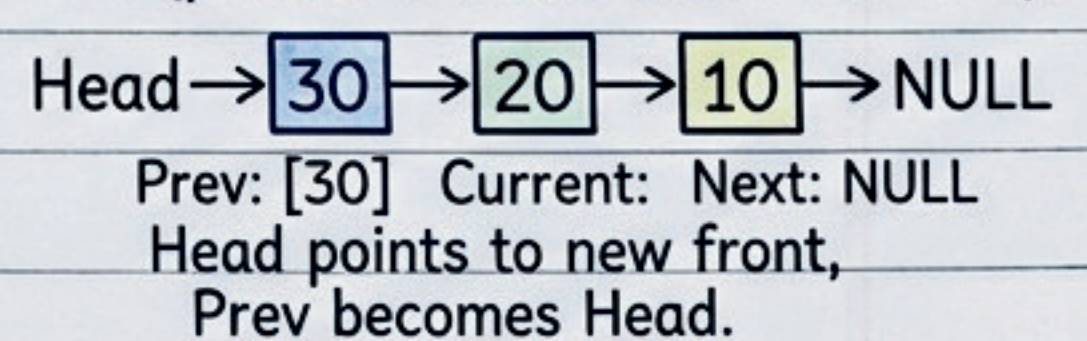
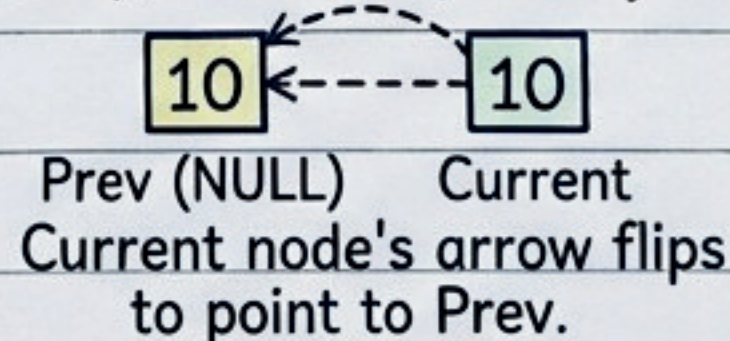
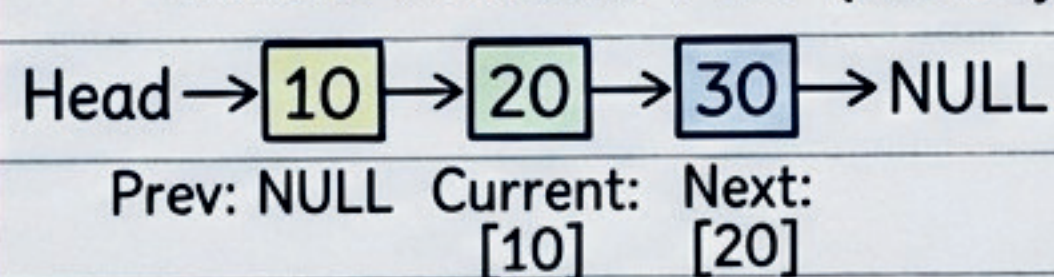
### Specific Position



## 3. Reversing a List

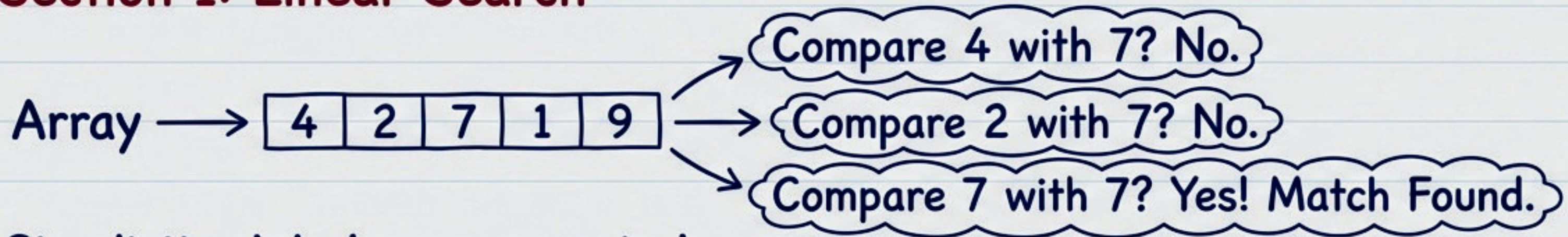
### How to reverse a Linked List?

- Define Pointers: Prev (initially NULL), Current (initially Head), Next (points to node after Current).



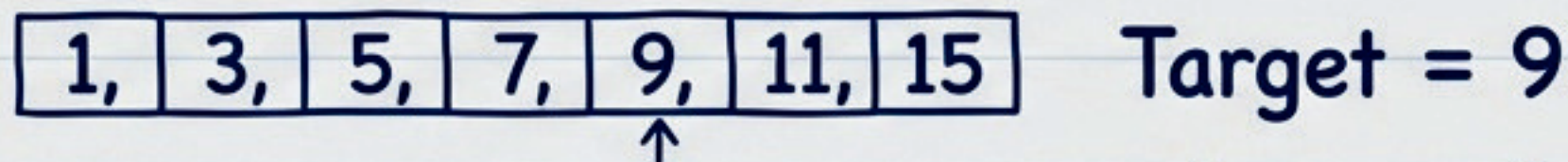
# Chapter 4: Searching Algorithms

## Section 1: Linear Search



Simplicity: Works on unsorted arrays.

## Section 2: Binary Search (Sorted Arrays Only)



Pass 1: Low=0, High=6. Mid=3 (Value 7).  $7 < 9$ , so "Search Right" →

Pass 2: New Low=4. New Mid (Value 11).  $11 > 9$ , so "Search Left" ←

Pass 3: Match 9! → 4

$Mid = \frac{Low + High}{2}$

## Section 3: Application (Example: Smallest Divisor)

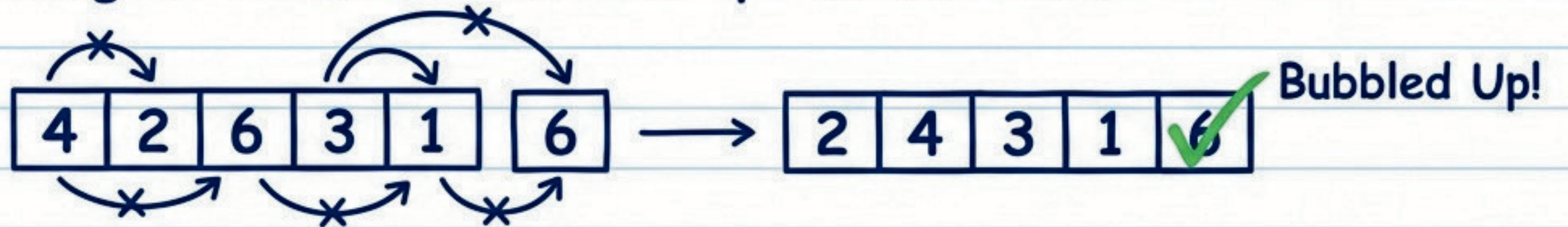
Find smallest divisor such that sum of quotients  $\leq 7$



# Chapter 5: Elementary Sorting Algorithms

## 1. Bubble Sort

Concept: Largest element "bubbles up" to the end.



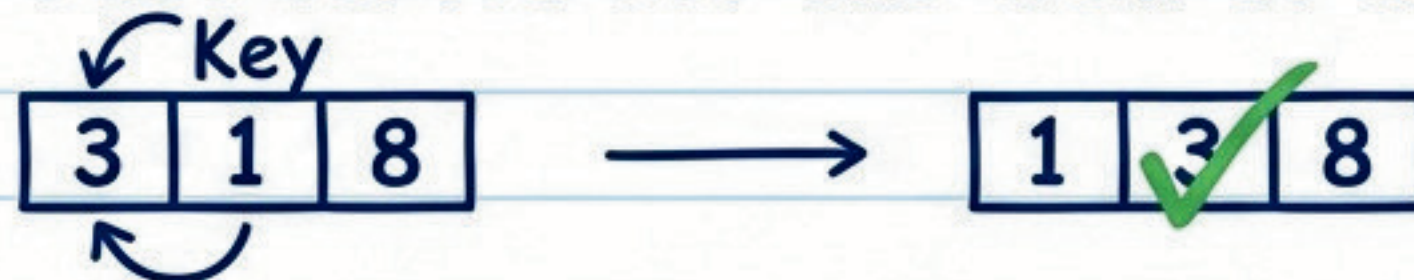
## 2. Selection Sort

Concept: Find minimum and swap to front.



## 3. Insertion Sort

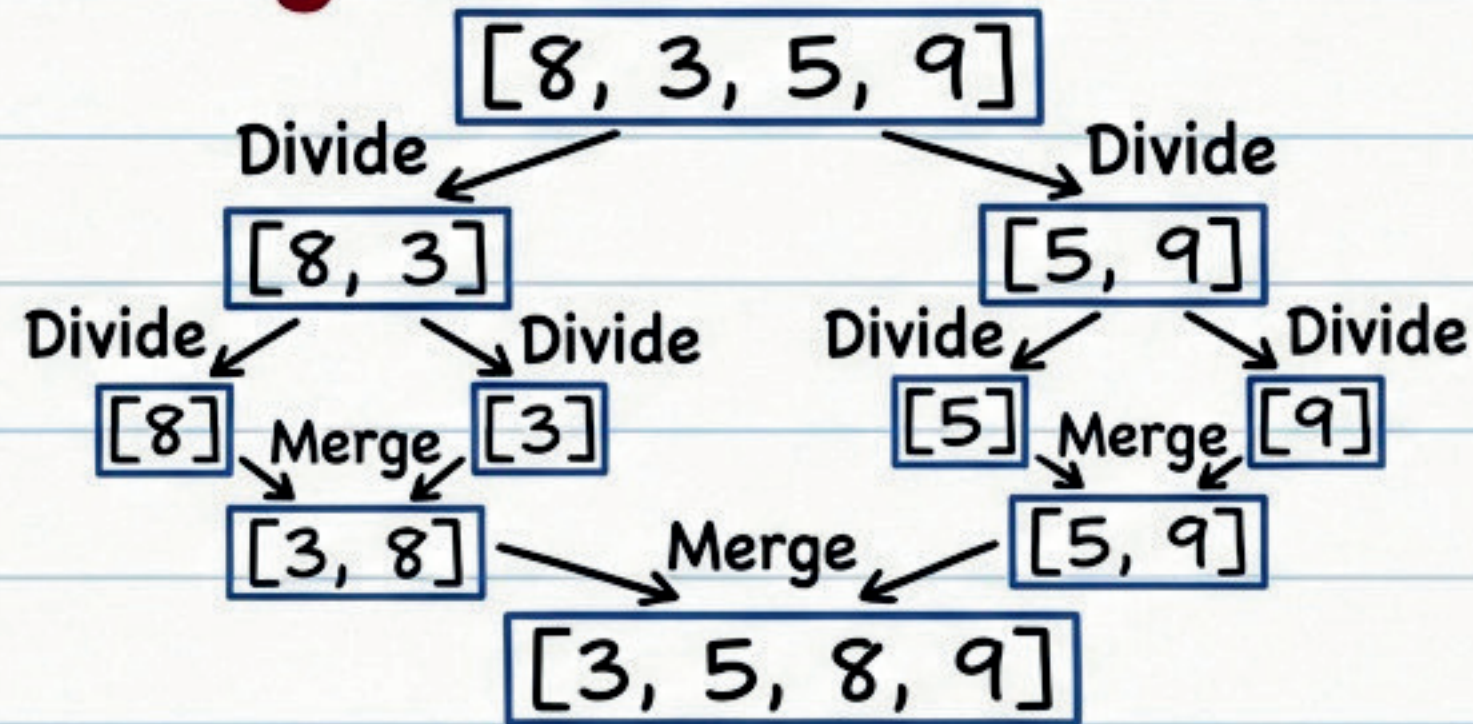
Concept: Build sorted list one item at a time (like cards).



All generally  $O(n^2)$ .

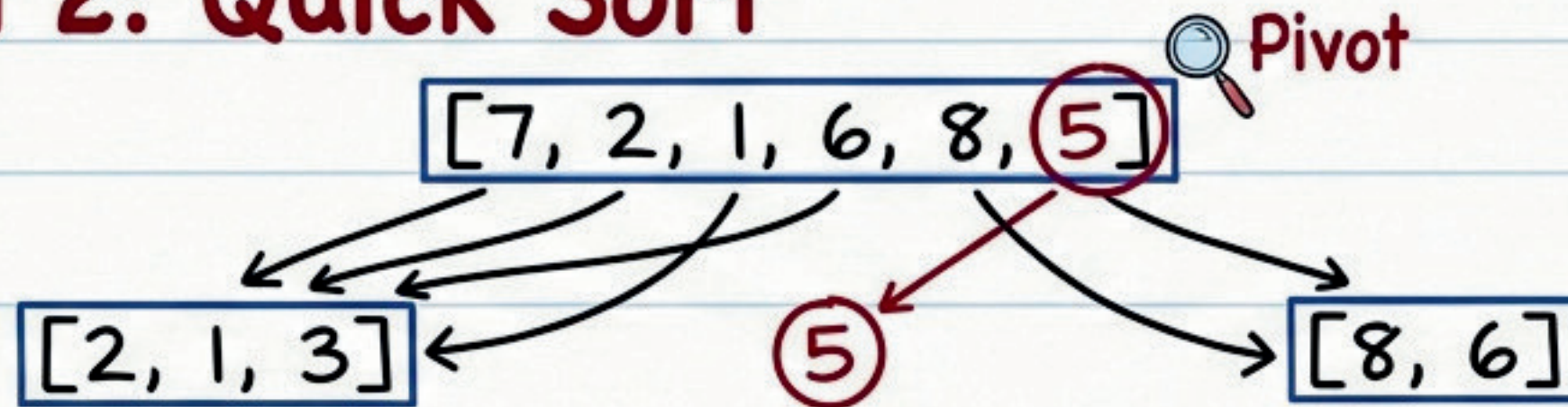
# Divide & Conquer Sorting

## Section 1: Merge Sort



**Stable**      **Time:  $O(n \log n)$**

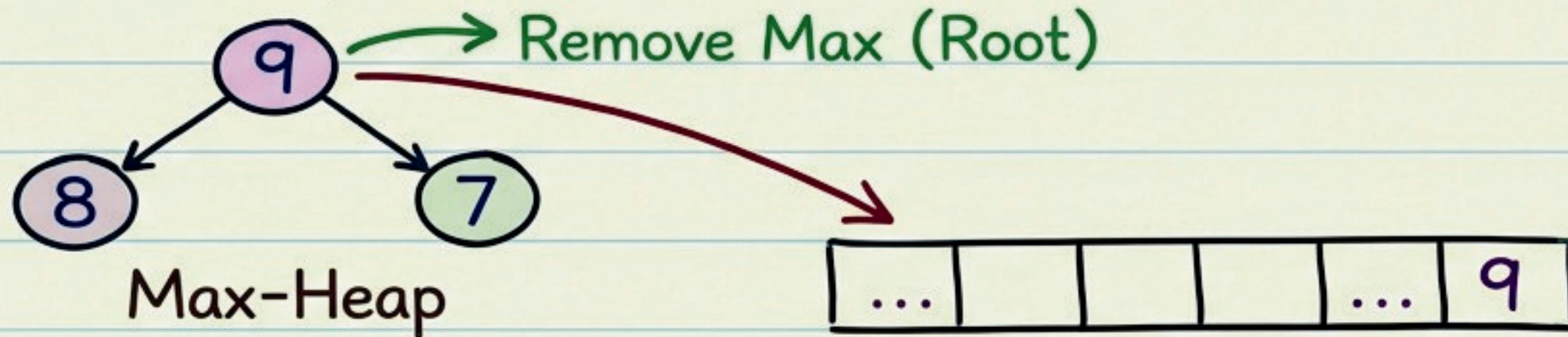
## Section 2: Quick Sort



**Not Stable**      **Avg Time:  $O(n \log n)$**       **Worst:  $O(n^2)$**

# Special Sorting Techniques

## Section 1: Heap Sort



💡 Algorithm: Build Max Heap → Swap Root with Last → Heapify

## Section 2: Counting Sort

Concept: Non-comparison based.

Input: [3, 5, 3, 1]

Count Array

1	2	3	4	5
1		2		1

Output: [1, 3, 3, 5]

✓ Complexity:  $O(n + k)$

# Radix Sort & Stability

## Section 1: Radix Sort

Sorts numbers digit by digit (LSD to MSD).

Step 1 (Units): Unsorted list: 

170	45	75	90	802
-----	----	----	----	-----

Sorted by last digit: 

170	90	802	45	75
-----	----	-----	----	----

Step 2 (Tens): Sorted by second digit.

(Tens): Sorted by second digit.  $[170, 45, 75, 90, 802] \rightarrow [45, 75, 170, 802, 90]$

Step 3 (Hundreds): Sorted by first digit.  $[45, 75, 170, 802, 90] \rightarrow [45, 75, 90, 170, 802]$

💡 Uses Counting Sort as a subroutine.

## Section 2: Stability

Definition: A sort is stable if it preserves the relative order of equal elements.

Unsorted: 

4A	2	4B
----	---	----

Stable Sort: 

2	4A	4B
---	----	----

(Correct)

Unstable Sort: 

2	4B	4A
---	----	----

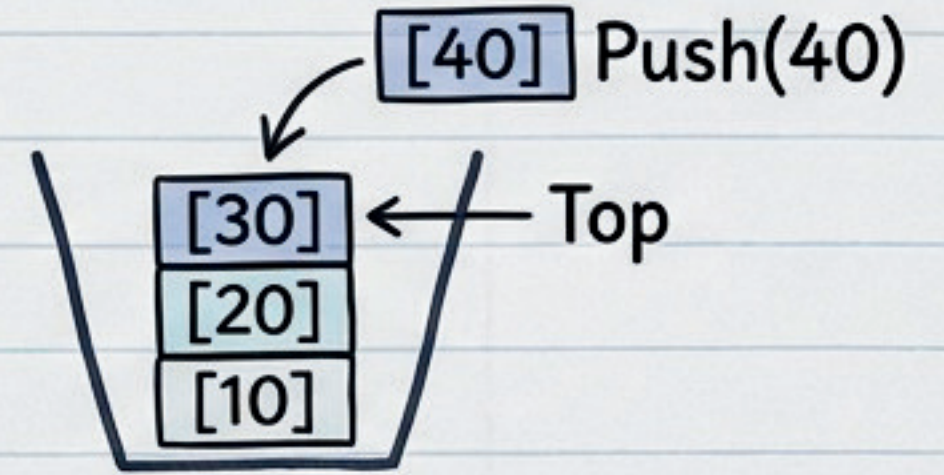
(Order swapped)

Radix Sort requires stability!

# Chapter 6: The Stack

## Section 1: Concept

✓ Definition: LIFO (Last In, First Out) principle.



## Section 2: Operations

✓ Push: Add to top.

✓ Pop: Remove from top.

History of Operations 

10	20	30	40
----	----	----	----

 →

Order of Removal 

40	30	20	10
----	----	----	----

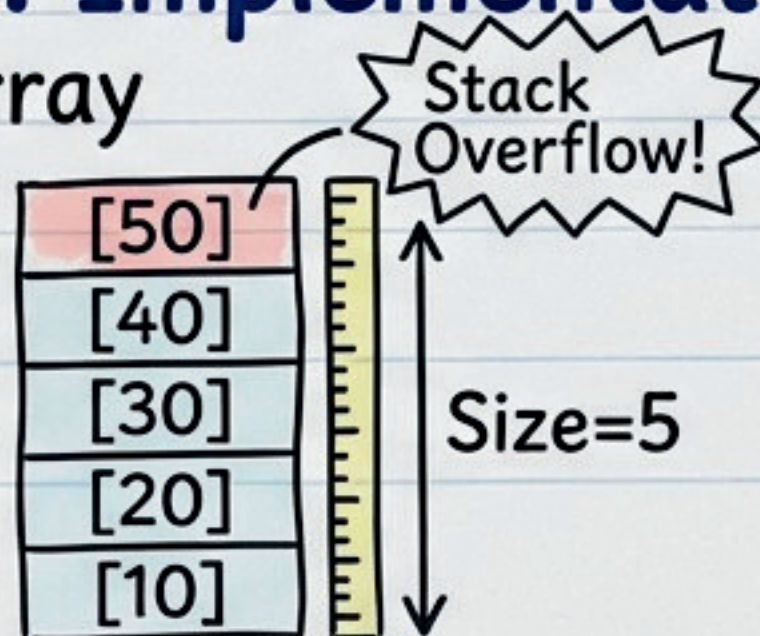
 →

LIFO Order: Last In (40), First Out

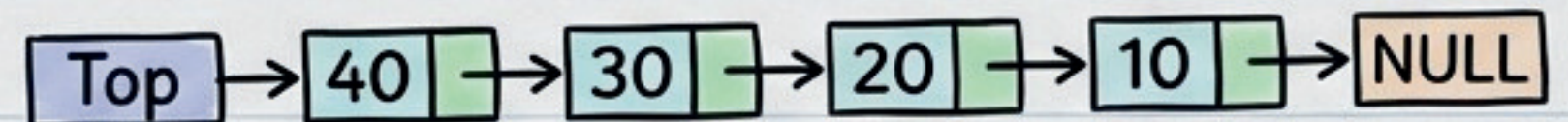
## Section 3: Implementation

Using Array

Stack  
(Fixed Size)



Using Linked List



Dynamic Size - No Overflow

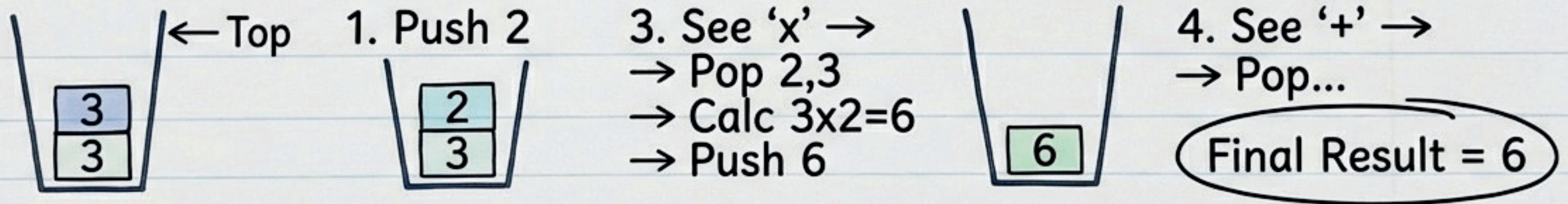
# Stack Applications & Expressions

## Section 1: Notations

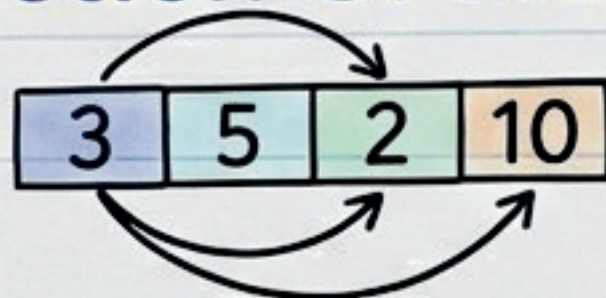
Notation	Operator Position	Example
Infix	Operator Between	$A + B$
Prefix	Operator Before	$+ A B$
Postfix	Operator After	$A B +$

## Section 2: Expression Evaluation

Evaluate Postfix using Stack:  $3\ 2\ x\ +$



## Section 3: Next Greater Element (NGE)

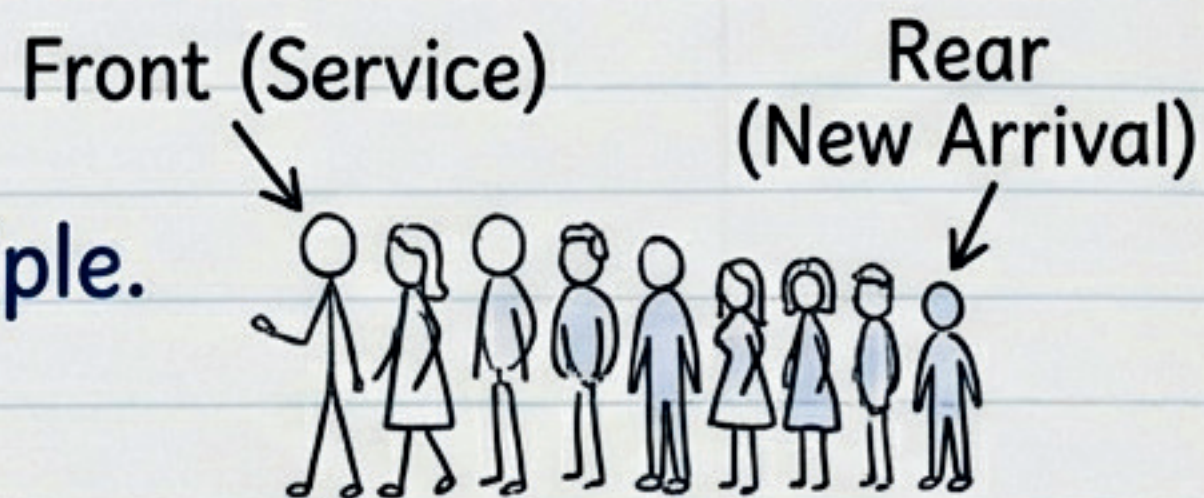


Find first greater element to the right using a Stack.

# Chapter 9: The Queue

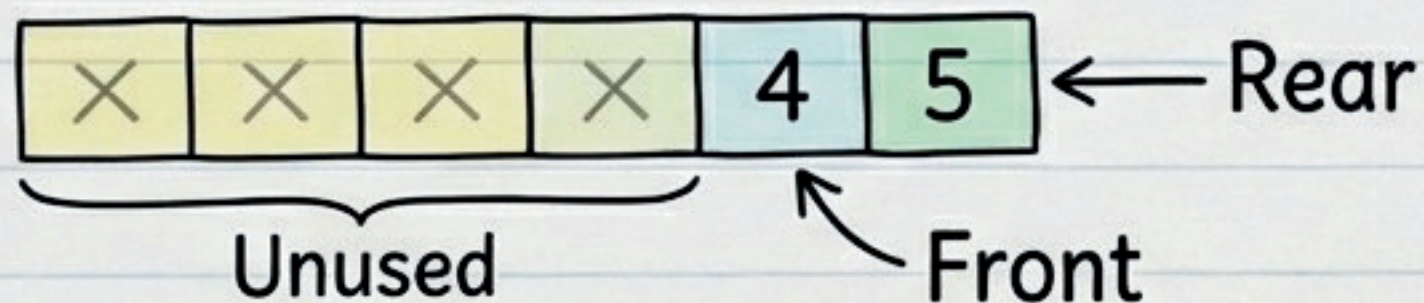
## Section 1: Concept

✓ Definition: FIFO (First In, First Out) principle.



## Section 2: The Problem

Simple Queue Limitation



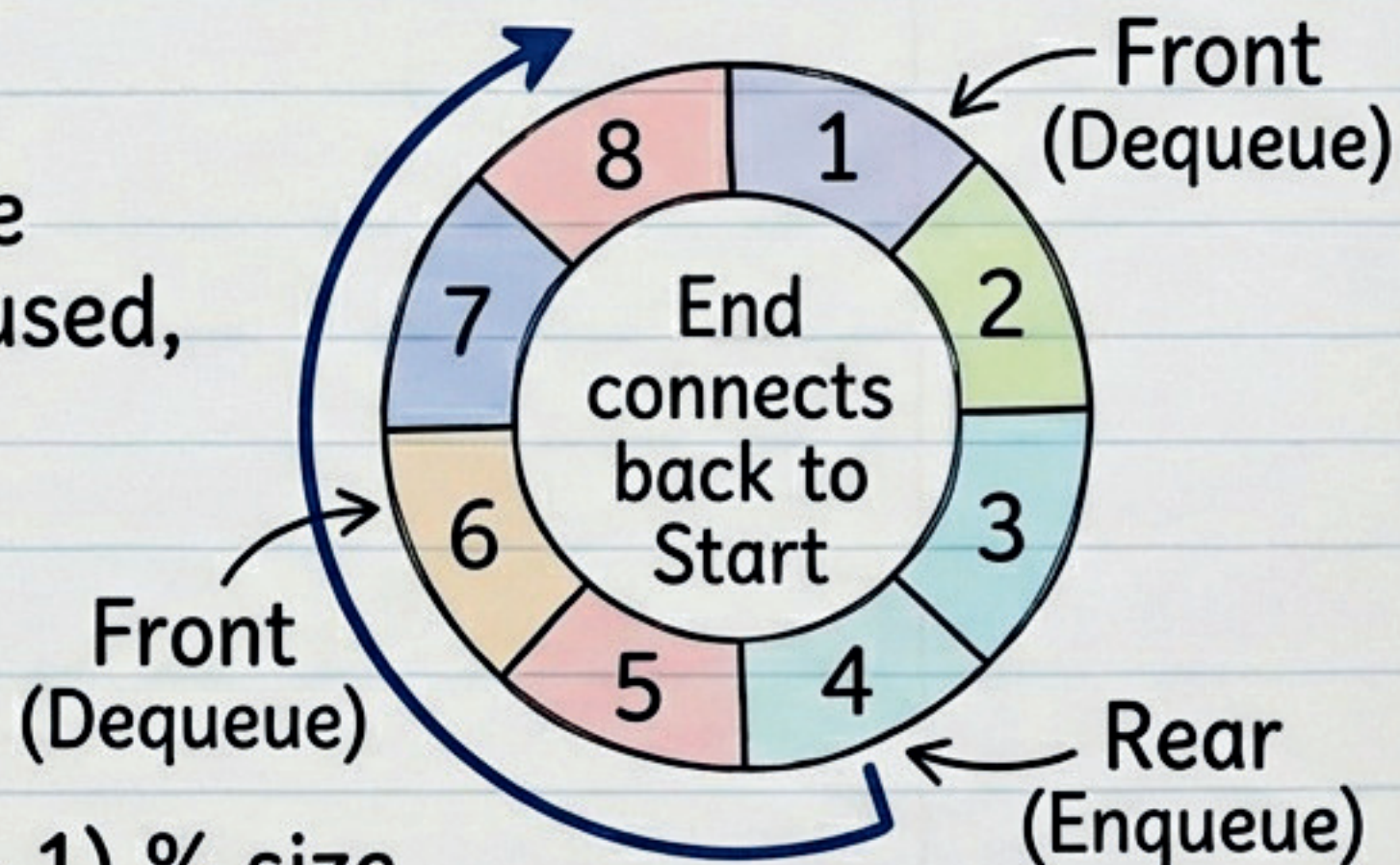
✓ Enqueue (Rear)  
✓ Dequeue (Front)

⚠ Once elements are removed from the front, the empty space cannot be reused, which leads to False Overflow.

## Section 3: The Solution

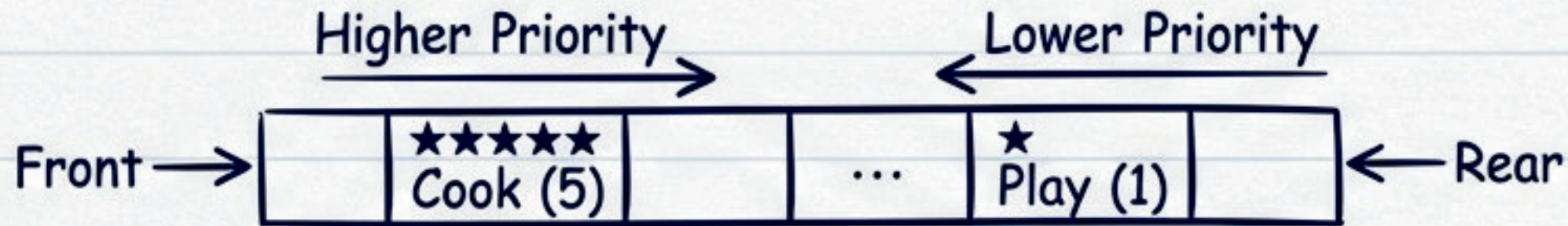
Circular Queue

Condition to move circularly:  $(\text{index} + 1) \% \text{size}$



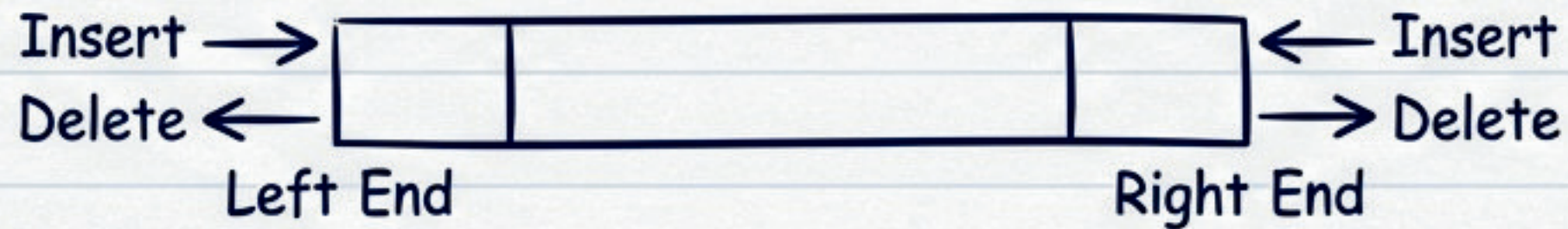
# Advanced Queues & Conversions

## Section 1: Priority Queue



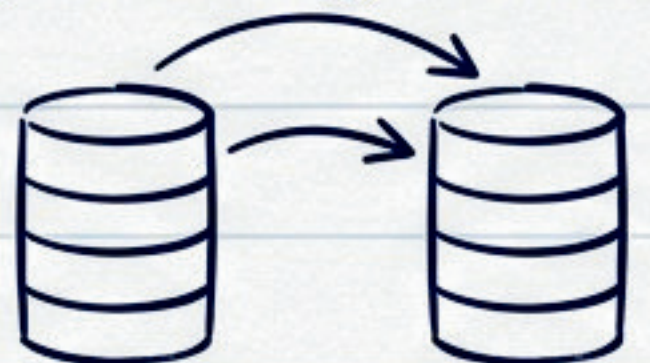
📌 Use Case: CPU Scheduling

## Section 2: Deque (Double Ended Queue)



## Section 3: Conversions

Queue using Stacks



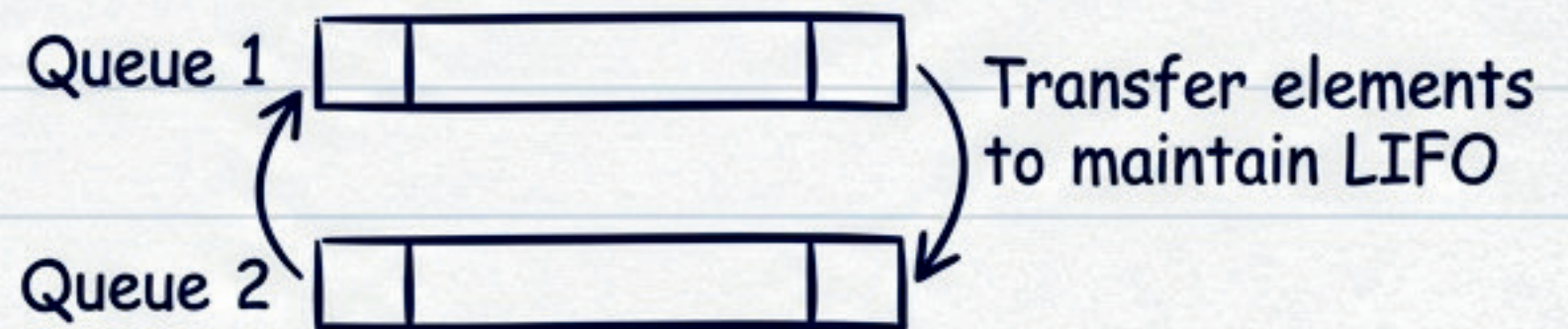
Stack 1

Stack 2

Stack 1 for enqueueing,  
Stack 2 for dequeueing

Reverse order by  
transferring

Stack using Queues Q1: Main, Q2: Auxiliary



Transfer elements  
to maintain LIFO